

What is claimed is:

1. A countermeasure method in an electronic component using a public key cryptography algorithm based on the use of elliptic curves in which a private key d and the number of points n on an elliptic curve are used to calculate a new deciphering integer d' such that the deciphering of any message, by means of a deciphering algorithm, with d' makes it possible to obtain the same result as with d , by performing the operation $Q=d.P$, where P is a point on the curve to which the scalar multiplication algorithm is applied, said method comprising the following steps:
 - 10 1) taking a random value r with the same size as d ;
 - 2) calculating an integer d' such that $d'=d+r$;
 - 3) Performing a scalar multiplication operation whose result is a point Q' on the curve such that $Q'=d'.P$;
 - 4) Performing a scalar multiplication operation whose result is a point S on the curve such that $S=r.P$; and
 - 15 5) calculating the point Q on the curve such that $Q=Q'-S$.
2. A countermeasure method according to Claim 1, wherein a new deciphering integer d' is calculated at each new execution of the deciphering algorithm.
- 20 3. A countermeasure method according to Claim 1, further including the step of incrementing a counter at each new execution of the deciphering algorithm up to an integer value T .
4. A countermeasure method according to Claim 3, wherein once the value T has been reached, a new deciphering integer d' is calculated

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according to the method of Claim 1, the counter is reset to zero and the point $S=r.P$ is stored in memory.

5. A countermeasure method according to Claim 4 wherein the value T is equal to 16.

5 6. A countermeasure method according to Claim 3 wherein the value T is equal to 16.

7. A countermeasure method according to Claim 1, wherein the point S is stored in memory, and steps 1 and 4 are replaced by the following steps 1' and 4':

- 10 1') replace r by 2.r
4') replace S by 2.S.

8. A countermeasure method according to Claim 7, wherein a new deciphering integer d' is calculated at each new execution of the deciphering algorithm.

15 9. A countermeasure method according to Claim 7, further including the step of incrementing a counter at each new execution of the deciphering algorithm up to a value T.

10. A countermeasure method according to Claim 9, wherein, once the value T has been reached, a new deciphering integer d' is calculated
20 according to the method of Claim 7, and the counter is reset to zero.

11. A countermeasure method according to Claim 10 wherein the value T is equal to 16.

12. A countermeasure method according to Claim 10 wherein the value T is equal to 16.

- 5 13. An electronic component having an integrated circuit which executes a public key cryptography algorithm based on the use of elliptic curves in which a private key d and the number of points n on an elliptic curve are used to calculate a new deciphering integer d' such that the deciphering of any message, by means of a deciphering algorithm, with d' makes it possible to
- 10 obtain the same result as with d, by performing the operation $Q=d.P$, where P is a point on the curve to which the scalar multiplication algorithm is applied, said circuit executing the following steps:
- 1) taking a random value r with the same size as d;
- 2) calculating an integer d' such that $d'=d+r$;
- 15 3) Performing a scalar multiplication operation whose result is a point Q' on the curve such that $Q'=d'.P$;
- 4) Performing a scalar multiplication operation whose result is a point S on the curve such that $S=r.P$; and
- 5) calculating the point Q on the curve such that $Q=Q'-S$.

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